## SECOND ORDER AND THIRD ORDER CONNECTIVITY INDICES OF A POLYPHENYLENE DENDRIMER

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The *m*-order connectivity index  ${}^{m}\chi(G)$  of a graph G is  $(d_{i_1}d_{i_2}...d_{i_{m+1}})^{-\frac{1}{2}}$ , where

 $d_{i_1}d_{i_2}\dots d_{i_{m+1}}$  runs over all paths of length *m* in *G* and  $d_i$  denotes the degree of vertex  $v_i$ .

A dendrimer is an artificially manufactured or synthesized molecule built up from branched units called monomers. In this paper, we compute 2- and 3-order connectivity index of an infinite family of polyphenylene dendrimer.

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## 1. Introduction

A simple graph G = (V,E) is a finite nonempty set V(G) of objects called vertices together with a (possibly empty) set E(G) of unordered pairs of distinct vertices of G called edges. In chemical graphs, the vertices of the graph correspond to the atoms of the molecule, and the edges represent the chemical bonds.

A single number which characterizes the graph of a molecular is called a graph theoretical invariant or topological index. Among the many topological indices considered in chemical graph, only a few have been found noteworthy in practical application, connectivity index is one of them. The connectivity index is one of the most popular molecular-graph. This index has been used in a wide spectrum of applications ranging from predicting physicochemical properties such as boiling point and solubility partition. The molecular connectivity index  $\chi$  provides a quantitative assessment of branching of molecules. Randic (1975) first addressed the problem of\_relating the physical properties of alkanes to the degree of branching across an isomeric series [6]. The degree of branching of a molecule was quantified using a branching index which subsequently became known as first- order molecular connectivity index  $\chi$ . Kier and Hall (1986) extended this to higher orders and introduced modifications to account for heteroatoms [4].

Molecular connectivity indices are the most popular class of indices (Trinajastic, 1992). They have been used in a wide spectrum of applications ranging from predicting physicochemical properties such as boiling point, solubility partition, coefficient etc, (Murray et al., 1975; Kier and Hall, 1976) for predicting biological activities such as antifungal effect, an esthetic effect, enzyme inhibition etc, (Kier et al., 1975; Kier and Murray, 1975) [4].

Let G be a simple connected graph of order n. For an integer  $m \ge 1$ , the m-order connectivity index of an organic molecule whose molecule graph G is defined as

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$$^{m}\chi(G) = \sum_{i_{1}...i_{m+1}} \frac{1}{\sqrt{d_{i_{1}}...d_{i_{m+1}}}}$$

where  $i_1...i_{m+1}$  runs over all paths of length *m* in *G* and  $d_i$  denote the degree of vertex  $v_i$ . In particular, 2-order connectivity index and 3-order connectivity index are defined as follows:

$${}^{2}\chi(G) = \sum_{i_{1}i_{2}i_{3}} \frac{1}{\sqrt{d_{i_{1}}d_{i_{2}}d_{i_{3}}}}, \quad {}^{3}\chi(G) = \sum_{i_{1}i_{2}i_{3}i_{4}} \frac{1}{\sqrt{d_{i_{1}}d_{i_{2}}d_{i_{3}}d_{i_{4}}}}.$$

Dendrimers are hyper-branched macromolecules, with a rigorously tailored architecture. They can be synthesized, in a controlled manner, either by a divergent or a convergent procedure. Dendrimers have gained a wide range of applications in supra-molecular chemistry, particularly in host guest reactions and self-assembly processes. Their applications in chemistry, biology and nano-science are unlimited. Recently, some researchers investigated *m*-order connectivity indices of some dendrimer nanostars, where m = 2 and 3 (see [1,2,3,7]). In this paper, we shall study the 2-and 3-order connectivity index of an infinite family of polyphenylene dendrimers.

## 2. Second-order and third-order connectivity index of dendrimer

In this section, we shall study the 2-order and 3-order connectivity index of a dendrimer. We consider polyphenylene dendrimer by construction of generations  $G_n$  with n growth stages. We denote this graph by  $D_4[n]$ . Figure 1 shows the generations  $G_2$  with 2 growth stages.



Fig. 1. Polyphenylene dendrimer of generations  $G_n$  with 2 growth stages.

The following theorem gives the 2-order connectivity index of polyphenylene dendrimer. **Theorem 1.** Let  $n \in \mathbb{N}$ . Then, the 2-order connectivity index of  $D_4[n]$  is given by

$${}^{2}\chi(D_{4}[n]) = \frac{1}{9}(198\sqrt{2} + 152\sqrt{3} + 6\sqrt{6} + 9) + (\frac{19\sqrt{2}}{2} + \frac{67\sqrt{3}}{9})(2^{n+1} - 4).$$

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**Proof.** First we compute  ${}^{2}\chi(D_{4}[1])$ . Let  $d_{i_{1}i_{2}i_{3}}$  denote the number of 2-paths whose three consecutive vertices are of degree  $i_1, i_2, i_3$ , respectively. In the same way, we use  $d_{i_1i_2i_3}^{(n)}$  to mean  $d_{i_1i_2i_3}$  in n-th stages. Particularly,  $d_{i_1i_2i_3}^{(n)} = d_{i_3i_2i_1}^{(n)}$ . We can see that

$$d_{222}^{(1)} = 48$$
,  $d_{223}^{(1)} = 48$ ,  $d_{232}^{(1)} = 24$ ,  $d_{233}^{(1)} = 56$ ,  $d_{323}^{(1)} = 4$ ,  $d_{333}^{(1)} = 44$ ,  $d_{234}^{(1)} = 8$ ,  $d_{343}^{(1)} = 6$ .

Therefore, we have

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$${}^{2}\chi(D_{4}[1]) = \frac{48}{\sqrt{2 \times 2 \times 2}} + \frac{48}{\sqrt{2 \times 2 \times 3}} + \frac{24}{\sqrt{2 \times 3 \times 2}} + \frac{56}{\sqrt{2 \times 3 \times 3}} + \frac{4}{\sqrt{3 \times 2 \times 3}} + \frac{44}{\sqrt{3 \times 3 \times 3}} + \frac{8}{\sqrt{2 \times 3 \times 4}} + \frac{6}{\sqrt{3 \times 4 \times 3}} = \frac{1}{9}(198\sqrt{2} + 152\sqrt{3} + 6\sqrt{6} + 9).$$

Now, we construct the relation between  ${}^{2}\chi(D_{4}[n])$  and  ${}^{2}\chi(D_{4}[n-1])$  for  $n \ge 2$ . By simple reduction, we have

$$\begin{aligned} d_{222}^{(n)} &= d_{222}^{(n-1)} + 18 \times 2^{n}, \ d_{223}^{(n)} &= d_{223}^{(n-1)} + 20 \times 2^{n}, \ d_{232}^{(n)} &= d_{232}^{(n-1)} + 10 \times 2^{n}, \ d_{233}^{(n)} &= d_{233}^{(n-1)} + 28 \times 2^{n}, \\ d_{323}^{(n)} &= d_{323}^{(n-1)} + 2 \times 2^{n}, \ d_{333}^{(n)} &= d_{333}^{(n-1)} + 22 \times 2^{n}, \end{aligned}$$

and for any  $(i_1i_2i_3) \neq (222), (223), (232), (233), (323), (333), (234), (343)$ , we have  $d_{i_1i_2i_3}^{(n)} = 0$ . Therefore

$${}^{2}\chi(D_{4}[n]) = {}^{2}\chi(D_{4}[n-1]) + \frac{18 \times 2^{n}}{\sqrt{2 \times 2 \times 2}} + \frac{20 \times 2^{n}}{\sqrt{2 \times 2 \times 3}} + \frac{10 \times 2^{n}}{\sqrt{2 \times 3 \times 2}} + \frac{28 \times 2^{n}}{\sqrt{2 \times 3 \times 3}} + \frac{2 \times 2^{n}}{\sqrt{3 \times 2 \times 3}} + \frac{22 \times 2^{n}}{\sqrt{3 \times 3 \times 3}}$$
$$= {}^{2}\chi(D_{4}[n-1]) + (\frac{19\sqrt{2}}{2} + \frac{67\sqrt{3}}{9}) \times 2^{n}.$$

From above recursion formula, we have

$${}^{2}\chi(D_{4}[n]) = {}^{2}\chi(D_{4}[n-1]) + (\frac{19\sqrt{2}}{2} + \frac{67\sqrt{3}}{9}) \times 2^{n}$$
  
=  ${}^{2}\chi(D_{4}[n-2]) + (\frac{19\sqrt{2}}{2} + \frac{67\sqrt{3}}{9})(2^{n} + 2^{n-1})$   
:  
=  ${}^{2}\chi(D_{4}[1]) + (\frac{19\sqrt{2}}{2} + \frac{67\sqrt{3}}{9})(2^{n} + 2^{n-1} + ... + 2^{2})$   
 ${}^{2}\chi(D_{4}[n]) = \frac{1}{9}(198\sqrt{2} + 152\sqrt{3} + 6\sqrt{6} + 9) + (\frac{19\sqrt{2}}{2} + \frac{67\sqrt{3}}{9})(2^{n+1} - 4).$ 

The proof is now complete.

The following theorem gives the 3-order connectivity index of polyphenylene dendrimer.

**Theorem 2.** Let  $n \in \mathbb{N}$ . Then, the 3-order connectivity index of  $D_4[n]$  is given by

$${}^{3}\chi(D_{4}[n]) = \frac{1}{9}(216 + 104\sqrt{6} + 6\sqrt{3} + 18\sqrt{2}) + \frac{1}{9}(99 + 46\sqrt{6})(2^{n+1} - 4).$$

**Proof.** Let  $d_{i_1i_2i_3i_4}$  denote the number of 3-paths whose four consecutive vertices are of degree  $i_1, i_2, i_3, i_4$ , respectively. With the same way, we use  $d_{i_1i_2i_3i_4}^{(n)}$  to mean  $d_{i_1i_2i_3i_4}$  in n-th stages. It is clear that  $d_{i_1i_2i_3i_4}^{(n)} = d_{i_4i_3i_2i_1}^{(n)}$ .

Similar to Theorem 1, we first compute  ${}^{3}\chi(D_{4}[1])$ . We can see that

$$\begin{split} &d_{2222}^{(1)}=32\;,\;d_{2223}^{(1)}=32\;,\;d_{2232}^{(1)}=48\;,\;d_{2233}^{(1)}=40\;,\;d_{2332}^{(1)}=16\;,\;d_{2333}^{(1)}=72\;,\;d_{3233}^{(1)}=16\;,\;,\\ &d_{3223}^{(1)}=8\;,\\ &d_{3233}^{(1)}=48\;,\;d_{2234}^{(1)}=8\;,\;d_{2343}^{(1)}=24. \end{split}$$

Thus,

$${}^{3}\chi(G_{4}[1]) = \frac{32}{\sqrt{2 \times 2 \times 2 \times 2}} + \frac{32}{\sqrt{2 \times 2 \times 2 \times 3}} + \frac{48}{\sqrt{2 \times 2 \times 3 \times 2}} + \frac{40}{\sqrt{2 \times 2 \times 3 \times 3}} + \frac{16}{\sqrt{2 \times 2 \times 3 \times 3}} + \frac{16}{\sqrt{3 \times 2 \times 3 \times 3}} + \frac{16}{\sqrt{3 \times 2 \times 3 \times 3}} + \frac{8}{\sqrt{3 \times 2 \times 2 \times 3}} + \frac{48}{\sqrt{3 \times 3 \times 3 \times 3}} + \frac{8}{\sqrt{2 \times 2 \times 3 \times 4}} + \frac{24}{\sqrt{2 \times 3 \times 4 \times 3}} = \frac{1}{9}(216 + 104\sqrt{6} + 6\sqrt{3} + 18\sqrt{2}).$$

Now, we compute  ${}^{3}\chi(D_{4}[n])$ .

The relations between  $d_{i_1i_2i_3i_4}^{(n)}$  and  $d_{i_1i_2i_3i_4}^{(n-1)}$  for  $n \ge 2$  are  $d_{i_1i_2i_3i_4}^{(n)} = d_{i_1i_2i_3i_4}^{(n-1)} + 12 \times 2^n - d_{i_1i_2i_3i_4}^{(n-1)} + 12 \times 2^n - d_{i_1i_2i_3i_4}^{(n-1)} + 20 \times 2^n$ 

$$d_{2222}^{(n)} = d_{2222}^{(n)} + 12 \times 2^{n}, \ d_{2223}^{(n)} = d_{2223}^{(n)} + 12 \times 2^{n}, \ d_{2232}^{(n)} = d_{2232}^{(n)} + 20 \times 2^{n},$$
  

$$d_{2233}^{(n)} = d_{2233}^{(n-1)} + 20 \times 2^{n},$$
  

$$d_{2332}^{(n)} = d_{2332}^{(n-1)} + 8 \times 2^{n}, \ d_{2333}^{(n)} = d_{2333}^{(n-1)} + 36 \times 2^{n}, \ d_{3223}^{(n)} = d_{3233}^{(n-1)} + 8 \times 2^{n}, \ d_{3223}^{(n)} = d_{3233}^{(n-1)} + 4 \times 2^{n},$$
  

$$d_{3333}^{(n)} = d_{3333}^{(n-1)} + 24 \times 2^{n},$$
  
and for any (ii i i i)  $\neq (2222)$  (2223) (2233) (2233) (2333) (3233)

and for any  $(i_1i_2i_3i_4) \neq (2222), (2223), (2232), (2233), (2332), (2333), (3233), (3223), (3333), (2234), (2343), we have <math>d_{i_1i_2i_3i_4}^{(n)} = 0$ . Therefore,

$${}^{3}\chi(D_{4}[n]) = {}^{3}\chi(D_{4}[n-1]) + \frac{12 \times 2^{n}}{\sqrt{2 \times 2 \times 2 \times 2}} + \frac{12 \times 2^{n}}{\sqrt{2 \times 2 \times 2 \times 3}} + \frac{20 \times 2^{n}}{\sqrt{2 \times 2 \times 3 \times 2}}$$

$$+ \frac{20 \times 2^{n}}{\sqrt{2 \times 2 \times 3 \times 3}} + \frac{8 \times 2^{n}}{\sqrt{2 \times 3 \times 3 \times 2}} + \frac{36 \times 2^{n}}{\sqrt{2 \times 3 \times 3 \times 3}} + \frac{8 \times 2^{n}}{\sqrt{3 \times 2 \times 3 \times 3}} + \frac{4 \times 2^{n}}{\sqrt{3 \times 2 \times 2 \times 3}} + \frac{24 \times 2^{n}}{\sqrt{3 \times 3 \times 3 \times 3 \times 3}}$$

$$= {}^{3}\chi(D_{4}[n-1]) + \frac{1}{9}(99 + 46\sqrt{6}) \times 2^{n}$$

$$= {}^{3}\chi(D_{4}[n-2]) + \frac{1}{9}(99 + 44\sqrt{6})(2^{n} + 2^{n-1})$$

$$\vdots$$

$$= {}^{3}\chi(D_{4}[1]) + \frac{1}{9}(99 + 44\sqrt{6})(2^{n} + 2^{n-1}... + 2^{2})$$

So,

$${}^{3}\chi(D_{4}[n]) = \frac{1}{9}(216 + 104\sqrt{6} + 6\sqrt{3} + 18\sqrt{2}) + \frac{1}{9}(99 + 46\sqrt{6})(2^{n+1} - 4).$$

The proof is now complete.

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